

Exercise 1 (Confluence & Commutation)

Show: If \rightarrow_1 and \rightarrow_2 are confluent, and if \rightarrow_1^* and \rightarrow_2^* commute, then $\rightarrow_{12} := \rightarrow_1 \cup \rightarrow_2$ is also confluent.

Exercise 2 (Strong Confluence)

A relation \rightarrow is said to be *strongly confluent* iff:

$$t_2 \leftarrow s \rightarrow t_1 \implies \exists u. t_2 \rightarrow^* u \leftarrow^* t_1$$

Show that every *strongly confluent* relation is also *confluent*.

Exercise 3 (Normal Forms)

Recall the inductive definition of the set NF of *normal forms*:

$$\frac{\frac{t \in \text{NF}}{\lambda x. t \in \text{NF}}}{n \geq 0 \quad t_1 \in \text{NF} \quad t_2 \in \text{NF} \quad \dots \quad t_n \in \text{NF}} \quad x \ t_1 \ t_2 \ \dots \ t_n \in \text{NF}$$

Show that this set precisely captures all normal forms, i.e.:

$$t \in \text{NF} \Leftrightarrow \nexists t'. t \rightarrow_\beta t'$$

Homework 4 (Semi-Confluence)

A relation \rightarrow is said to be *semi-confluent* iff:

$$t_2 \stackrel{*}{\leftarrow} s \rightarrow t_1 \implies \exists u. t_2 \rightarrow^* u \stackrel{*}{\leftarrow} t_1$$

Show that \rightarrow is *semi-confluent* if and only if it is *confluent*.

Homework 5 (Weak Diamond Property)

Assume that \rightarrow has the following weaker diamond property:

$$t_2 \leftarrow s \rightarrow t_1 \wedge t_1 \neq t_2 \implies \exists u. t_2 \rightarrow u \leftarrow t_1.$$

- a) Is it still the case that every element is either in normal form or has no normal form?
- b) Show that if t has a normal form, then all its reductions to its normal form have the same length.

Homework 6 (Normal Forms & Big Step)

Show:

$$t \in \text{NF} \wedge t \Rightarrow_n u \implies u = t$$