

Exercise 1 (Confluence & Commutation)

Show: If \rightarrow_1 and \rightarrow_2 are confluent, and if \rightarrow_1^* and \rightarrow_2^* commute, then $\rightarrow_{12} := \rightarrow_1 \cup \rightarrow_2$ is also confluent.

Exercise 2 (Strong Confluence)

A relation \rightarrow is said to be *strongly confluent* iff:

$$t_2 \leftarrow s \rightarrow t_1 \implies \exists u. t_2 \rightarrow^* u \leftarrow^* t_1$$

Show that every *strongly confluent* relation is also *confluent*.

Exercise 3 (Local Confluence of η -reduction)

Analogously to β -reduction, we define η -reduction inductively:

1. $x \notin \text{FV}(s) \implies (\lambda x. s x) \rightarrow_\eta s$
2. $s \rightarrow_\eta s' \implies s t \rightarrow_\eta s' t$
3. $s \rightarrow_\eta s' \implies t s \rightarrow_\eta t s'$
4. $s \rightarrow_\eta s' \implies (\lambda x. s) \rightarrow_\eta (\lambda x. s')$

The proof of local confluence of \rightarrow_η , i.e. it holds that there exists a u with $t_1 \rightarrow_\eta^* u \leftarrow_\eta^* t_2$ if we have $t_1 \leftarrow_\eta s \rightarrow_\eta t_2$, was very informal. Give a proper proof using this definition.

Homework 4 (Semi-Confluence)

A relation \rightarrow is said to be *semi-confluent* iff:

$$t_2 \xrightarrow{*} s \rightarrow t_1 \implies \exists u. t_2 \xrightarrow{*} u \xrightarrow{*} t_1$$

Show that \rightarrow is *semi-confluent* if and only if it is *confluent*.

Homework 5 (Diamond Property & Normal Forms)

Show that if \rightarrow has the diamond property, every element is either in normal form or has no normal form.

Homework 6 (Weak Diamond Property)

Assume that \rightarrow has the following weaker diamond property:

$$t_2 \leftarrow s \rightarrow t_1 \wedge t_1 \neq t_2 \implies \exists u. t_2 \rightarrow u \leftarrow t_1.$$

- a) Is it still the case that every element is either in normal form or has no normal form?
- b) Show that if t has a normal form, then all its reductions to its normal form have the same length.